

Irreducibility of configurations

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Incidence geometry

An **incidence geometry** is a triple (P, L, I) where

- P is a set of 'points',
- L is a set of 'blocks',
- I is an incidence relation between the elements in P and L .

When there are at most one block containing p_i and p_j for all pairs of points, then we call the blocks **lines**.

Incidence graph of an incidence geometry

We can use a graph to represent the incidences of points and blocks.

The **incidence graph** of the incidence structure (P, L, I) is the bipartite graph with vertex set $P \cup L$ and an edge between the vertices p and b if p is a point on b .

Combinatorial configurations

A **combinatorial** (v, b, r, k) -**configuration** is an incidence structure with v points and b lines/blocks such that

- every point appears on r lines,
- every line has k points,
- every pair of points is in at most one line, or equivalently,
- every pair of lines intersect in at most one point.

The four parameters (v, b, r, k) are redundant.

We only need the three parameters (d, r, k) with

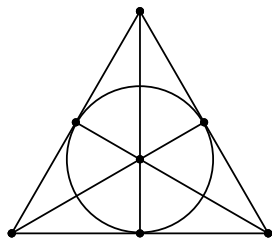
$$d := \frac{v \gcd(r, k)}{k} = \frac{b \gcd(r, k)}{r} = \frac{vr}{\text{lcm}(r, k)} = \frac{bk}{\text{lcm}(r, k)}.$$

Reduced parameters: (d, r, k) -configuration.

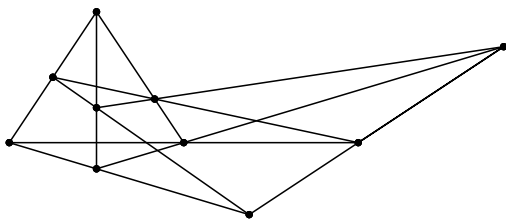
A combinatorial (v, b, r, k) configuration is also called an r -regular and k -uniform **partial linear space**.

Balanced configurations

We say that a combinatorial configuration is **balanced** if $r = k$. This implies that the number of points equals the number of lines and also, the associated integer, so $d = v = b$.



The Fano plane,
 $(v, b, r, k) = (7, 7, 3, 3)$
 $(d, r, k) = (7, 3, 3)$

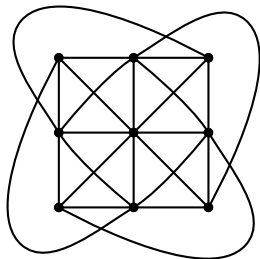


The Desargues' configuration
 $(v, b, r, k) = (10, 10, 3, 3)$
 $(d, r, k) = (10, 3, 3)$

Non-balanced configurations

When $r \neq k$, then $v \neq b$ and $d = \frac{v \gcd(r,k)}{k}$.

- The affine plane $AG(2, q)$ over the finite field \mathbb{F}_q has parameters $(q^2, q^2 + q, q + 1, q)$ so $d = q$.
Reduced parameters: $(q, q + 1, q)$.
- A Steiner triple systems of order v ($STS(v)$) has parameters $(v, v(v - 1)/6, (v - 1)/2, 3)$.
Reduced parameters: $(v \gcd(v - 1, 3)/3, (v - 1)/2, 3)$.



$AG(2, 3) / STS(9)$

$$(v, b, r, k) = (9, 12, 4, 3)$$

$$(d, r, k) = (3, 4, 3)$$

Necessary conditions for existence of configurations

The following necessary conditions for existence of configurations are well-known.

Lemma.

Suppose that there exists a (v, b, r, k) -configuration. Then

- ① $v \geq r(k - 1) + 1$ and $b \geq k(r - 1) + 1$, and
- ② $vr = bk$.

We say that parameters satisfying these conditions are **admissible**.

What about sufficient conditions?

Sufficient conditions

- When $r = 3$, the necessary conditions are sufficient [Gropp (1994)].
- When $r = 4$, it is conjectured that the necessary conditions are sufficient [Gropp (2001)].
- When $r = 5$, the necessary conditions are not sufficient. Sufficient conditions are not known for $k > r$.
- In general sufficient conditions are not known.

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Augmenting balanced configurations [Martinetti, 1886]

Given a $(v, v, 3, 3)$ -configuration, add a point and a line to construct a $(v + 1, v + 1, 3, 3)$ -configuration.

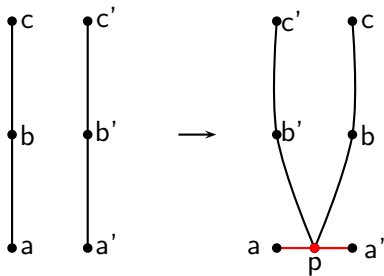
How?

Assume that there are two parallel lines $\{a, b, c\}$ and $\{a', b', c'\}$, with a and a' noncollinear.

Add a point p and replace the two parallel lines with the lines $\{p, b, c\}$, $\{p, b', c'\}$, $\{p, a, a'\}$.

The result is a $(v + 1, v + 1, 3, 3)$ -configuration.

The Martinetti augmentation



Reduction of configurations I [Martinetti, 1886]

A configuration is called irreducible if it cannot be constructed from a smaller configuration using the augmentation construction.

Theorem. [Martinetti- Boben]

The irreducible configurations à la Martinetti are:

- Cyclic configurations with base line $\{0, 1, 3\}$ (starting with the Fano plane).
- Three infinite families $T_1(n)$, $T_2(n)$, $T_3(n)$, on $10n$ points. The smallest configuration in $T_1(n)$ is the Desargues' configuration.
- The Pappus' configuration.

Reduction of configurations II [Carstens et al., 2001]

Given a $(v, v, 3, 3)$ -configuration, remove a point and a line to construct a $(v - 1, v - 1, 3, 3)$ -configuration.

How?

A complicated family of several Martinetti-like reductions defined on the incidence graph.

Their goal was to show that the only irreducible configuration was the Fano plane.

Unfortunately, in 2005, Ravnik used a computer to show that they failed to reduce at least the Desargues' configuration.

Reduction of configurations III [Boben, 2005]

In the incidence graph of a $(v, v, 3, 3)$ -configuration, remove a point-vertex p and a line-vertex ℓ and connect their neighbors so that the result is an incidence graph of a $(v - 1, v - 1, 3, 3)$ -configuration.

The incidence graph of a $(v, v, 3, 3)$ -configuration is a bipartite cubic graph of girth at least 6. The result is a $(v - 1, v - 1, 3, 3)$ -configuration.

Martinetti's reduction is a special case of this reduction.

Irreducible configurations

Theorem. [Boben (2005)] Boben's irreducible configurations are:

- The Fano plane.
- The Pappus' configuration.

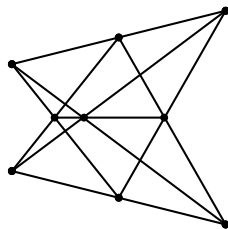
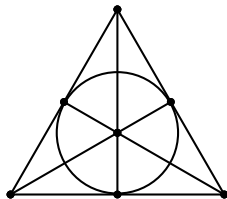


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Augmenting balanced configurations with $r = k \geq 3$

Theorem.

Assume there is

- a (balanced) (d, k, k) -configuration (P, L, I) with k points $Q \subseteq P$ and k lines $M \subseteq L$, and
- a bijection $f : Q \rightarrow M$ defined as follows:
 - ▶ the image of a point $q \in Q$ is a line $f(q) \in M$ through that point,
 - ▶ two points $q, q' \in Q$ can be collinear only on the line $f(q)$ or $f(q')$,
 - ▶ two lines $m, m' \in M$ can meet only in the point $f^{-1}(m)$ or $f^{-1}(m')$.

Then there is a $(d + 1, k, k)$ -configuration constructed from C through an augmentation procedure.

Augmenting balanced configurations with $r = k \geq 3$

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 - ▶ two lines $m, m' \in M$ can meet only in the point $f^{-1}(m)$ or $f^{-1}(m')$.

Then there is a $(d + 1, k, k)$ -configuration constructed from C through an augmentation procedure.

- Disconnect all incidences $(q, f(q)) \in I$.
- Add a new line ℓ and the incidences (q, ℓ) for all points $q \in Q$.
- Add a new point p and the incidences (p, m) for all lines $m \in M$.

The result is a configuration with parameters $(d + 1, k, k)$.

Example: Augmenting $(v, v, 3, 3)$ -configurations

Lemma. Every $(v, v, 3, 3)$ -configuration admits an augmentation.

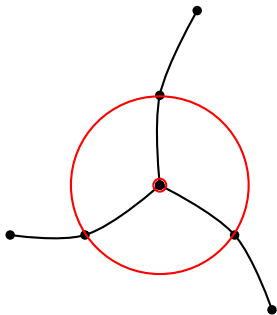
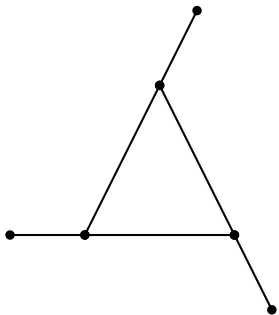
Proof.

- If the configuration contains a triangle, take Q and M the three points and the three lines in the triangle.
- If the configuration contains no triangle, there are still three points a, b, c such that (a, b) and (b, c) are collinear on the two lines A, B , and a third line C through a not meeting A nor B . □

This implies the following well-known result.

Corollary. There is a $(v, v, 3, 3)$ configuration whenever the parameters are admissible.

Augmenting $(v, v, 3, 3)$ -configurations



Deficiency of a configuration

The **distance** between two points is the number of lines in a “shortest path” between them.

In a (d, r, k) -configuration all points have the same number $r(k - 1)$ of points at distance 1.

The **deficiency** of a configuration is the number of points at distance at least 2 from a given point.

Augmenting $(v, v, 4, 4)$ -configurations

Lemma. A $(d, 4, 4)$ -configuration admits an augmentation if and only if it has deficiency at least 1.

Proof.

- If deficiency is 0 then it is the finite projective plane of order 3, which is not augmentable.
- If deficiency is ≥ 1 then there are always points a, b, c, d such that the pairs (a, b) , (b, c) , (c, d) are collinear on the lines A, B, C , and the pairs (a, c) , (b, d) are at distance at least two. The fourth line D can be taken as the line through a and d if there is such a choice of points and lines. Otherwise D can be taken through d such that it does not meet A, B, C . □

There are $(v, 4, 4)$ -configuration with deficiency 0 and 1, so we get the following well-known result.

Corollary. There is a $(v, v, 4, 4)$ -configuration whenever the parameters are admissible.

Augmenting configurations with $r, k \geq 3$

Theorem. Let $t = rk / \gcd(r, k)$. Assume there is

- a (d, r, k) -configuration (P, L, I) with t points $Q \subseteq P$ and t lines $M \subseteq L$, and
- a bijection $f : Q \rightarrow M$ defined as follows:
 - ▶ the image of a point $q \in Q$ is a line $f(q) \in M$ through that point,
 - ▶ $Q = \bigcup_{i=1}^{r/\gcd(r,k)} Q_i$ such that $|Q_i| = k$, $Q_i \cap Q_j = \emptyset$, and two points $q_i, q'_i \in Q_i$ can be collinear only on the line $f(q_i)$ or $f(q'_i)$,
 - ▶ $M = \bigcup_{i=1}^{k/\gcd(r,k)} M_i$ such that $|M_i| = r$, $M_i \cap M_j = \emptyset$, and two lines $m_i, m'_i \in M_i$ can meet only in the point $f^{-1}(m_i)$ or $f^{-1}(m'_i)$.

Then there is a $(d + 1, r, k)$ -configuration constructed from C through an augmentation procedure.

Augmenting configurations with $r, k \geq 3$

Proof.

- Disconnect all incidences $(q, f(q)) \in I$.
- For each Q_i add a new line l_i and the incidences (q_i, l_i) for all points $q_i \in Q_i$.
- For each M_i add a new point p_i and the incidences (p_i, m_i) for all lines $m_i \in M_i$.

The result is a configuration with parameters $(d + 1, r, k)$.

$$\begin{aligned}(d, r, k) &\rightarrow (d + 1, r, k) \\ (v, b, r, k) &\rightarrow (v + k/\gcd(r, k), b + r/\gcd(r, k), r, k)\end{aligned}$$



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Reduction of **balanced** configurations with $r = k \geq 3$

A **reduction** of a balanced configuration (P, L, I) is a triple (p, ℓ, f') where

- p is a point,
- ℓ is a line,
- f' is a bijection $f' : Q' \rightarrow M'$,
where
 - ▶ $Q' = \{q : q \in \ell \text{ and } q \neq p\}$, and
 - ▶ $M' = \{m : p \in m \text{ and } m \neq \ell\}$,

such that q is not collinear with $s \in f'(q)$ except possibly through ℓ or with p .

Now delete p and I (and their incidences) and add incidences $(q, f'(q))$ for $q \in Q'$.

A balanced configuration is **irreducible** if it does not admit a reduction.

Lemma. The reduction is the inverse operation of the augmentation.

Reduction of configurations with $r, k \geq 3$

A **reduction** of a configuration (P, L, I) is a triple (R, N, f') where

- R is a set of points,
- N is a set of lines,
- f' is a pairing between the elements of two multisets $f : Q' \rightarrow M'$, where
 - ▶ $Q' = \{q : q \in P \text{ and } \exists \ell \in N \text{ such that } q \in \ell \text{ and } q \notin R\}$,
 - ▶ $M' = \{m : m \in L \text{ and } \exists p \in R \text{ such that } p \in m \text{ and } m \notin N\}$,

such that q is not collinear with $s \in f'(q)$ except possibly through one of the lines in N or with one of the points in R .

Now delete R and N and their incidences and add incidences $(q, f'(q))$ for $q \in Q'$.

A configuration is **irreducible** if it does not admit a reduction.

Lemma. The reduction is the inverse operation of the augmentation.

Reduction of $(d, 3, 3)$ -configurations

In the case $(d, 3, 3)$ this definition has the same implications as Boben's reduction.

Lemma There are only two irreducible $(d, 3, 3)$ -configurations:

- The Fano plane,
- The Pappus' configuration.

The Pappus' configuration as a transversal design

What is the analog of the Pappus' configuration for other values of r and k ?

The Pappus' configuration is a **resolvable transversal design**.

A transversal design $TD_\lambda(k, n)$ is a (k -uniform) incidence geometry on $v = kn$ points partitioned into k groups of n elements, such that

- any group and any block contain exactly one common point, and
- every pair of points from distinct groups is contained in exactly λ blocks.

A transversal design is **resolvable** if the line set can be partitioned in parallel classes and it is a (kn, n^2, n, k) -configuration if $\lambda = 1$.

Example: There is a resolvable $TD_1(k, n)$ whenever there is an affine plane of order n and $k \leq n$. Take the points on k lines in a parallel class and restrict the rest of the lines to these points.

Irreducibility of resolvable transversal designs

Lemma. A resolvable transversal design $TD_1(k, n)$ is irreducible if $k \geq (k + r)/\gcd(r, k) + 1$.

Proof.

- Let p be a point in $T = TD_1(k, n)$ and ℓ_1, \dots, ℓ_n the lines through p .
- Then ℓ_1, \dots, ℓ_n are in different parallel classes.
- Let ℓ be a line in T and q a point on ℓ .
- Then q is collinear with all points on the lines ℓ_1, \dots, ℓ_n except one on each line.
- At most $(r + k)/\gcd(r, k)$ of these incidences can be removed by a reduction and do (perhaps) not obstruct reduction.
- More than $(r + k)/\gcd(r, k)$ such incidences will obstruct reduction.

Large configurations are reducible

Lemma. A (v, b, r, k) -configuration is reducible if $b \geq 1 + r + r(k - 1)(r - 1) + r(k - 1)^2(r - 1)^2$.

However this bound is not sharp.

Thank you for listening!